

On distinguishing proof-theoretic consequence from derivability

Nissim Francez,

Computer Science dept., the Technion-IIT, Haifa, Israel (francez@cs.technion.ac.il)

According to the common conception of *logical consequence*, it can be defined in two main ways:

- **Model-theoretically:** For a suitable notion of a model, consequence is taken to be *preservation* (also called *propagation or transmission*) of truth over all models.

$$\Gamma \models \varphi \text{ iff for every model } \mathcal{M}, \text{ if } \mathcal{M} \models \Gamma \text{ then } \mathcal{M} \models \varphi \quad (1)$$

- **Proof-Theoretically:** For a suitable complete proof-system \mathcal{N} , consequence is taken as *derivability* in \mathcal{N} , denoted ' $\vdash_{\mathcal{N}}$ '.

$$\Gamma \vdash_{\mathcal{N}} \varphi \text{ iff there exists a derivation } \mathcal{D} \text{ of } \varphi \text{ from (open) assumptions } \Gamma \quad (2)$$

The idea that logical consequence involves preservation of *something*, not necessarily of truth, has been suggested by many. Some examples:

- *Information:* Propagation of the information (in a situation) is underlying consequence of Relevant Logic.

- *Ambiguity:* a proposition p is taken as *ambiguously* between two different propositions, p_t and p_f . A *measure* of ambiguity of an inconsistent Γ is defined as the minimal number of proposition in Γ the treatment of which as ambiguous renders Γ consistent. Propagation of ambiguity is used for defining consequence for paraconsistent logics.

- *Precisification:* In the context of *vagueness*, Logical consequence is preservation of *super-truth* (i.e., truth in all *precisifications*).

A natural question arising is, what is common to all the “things” being suggested as preserved, or propagated, by the various consequence relations mentioned above?

I want to argue that they all serve (either explicitly or implicitly) as *central concepts on which theories of meaning are based*.

Two of the main theories of meaning are the following.

– In **Model-Theoretic Semantics** (MTS), The central concept is *truth* (in arbitrary models). Meaning is defined as truth-conditions.

– In **Proof-Theoretic Semantics** (PTS), The central concept is *proof*, or more generally, *canonical derivation*, in appropriate meaning-conferring proof-systems.:

Meaning is *determined* (either implicitly or, as I prefer, explicitly) by the rules of the meaning-conferring system.

– The other propagated “things” mentioned above have a similar role as theories of meaning for Relevant Logic, general paraconsistent logics and for languages with vagueness.

Consequently, I suggest the following informal principle:

meaning-based logical consequence: In a theory of meaning \mathcal{T} , logical consequence is based on the propagation of the central concept of \mathcal{T} .

In this paper I argue that, in spite of the coextensiveness in many logics of derivability and preservation of truth in models, if one adheres to the proof-theoretic semantics theory of meaning then (2) is not the right definition of proof-theoretic consequence. While (1) is faithful to the usual model-theoretic conception of meaning, (2) is not faithful to the PTS conception of meaning.

Instead, I suggest the following conception of proof-theoretic consequence.

– The (*proof-theoretic*) *meaning* $\llbracket \varphi \rrbracket$ of φ , is given by: $\llbracket \varphi \rrbracket = \lambda \Gamma. \llbracket \varphi \rrbracket_{\Gamma}^c$, a function from contexts Γ to the (possibly empty) collection of *canonical derivations* of φ from Γ ($\Gamma \vdash^c \varphi$).

– The definition of *proof-theoretic consequence* (pt-consequence) rests on the notion of *grounds for assertion* for φ , closely related to $\llbracket \varphi \rrbracket$, given by: $GA[\varphi] = \{\Gamma \mid \Gamma \vdash^c \varphi\}$.

– Proof-theoretic consequence: ψ is a *proof-theoretic consequence* of Γ ($\Gamma \Vdash \psi$) iff $GA[\Gamma] \subseteq GA[\psi]$.

The paper studies two definitions of $GA[\Gamma]$, based on conjunction (additive) and on fusion (multiplicative).

I show that for intuitionistic logic, but not for classical logic, proof-theoretic consequence coincides with derivability.